



COLLÈGE
DE FRANCE
—1530—



Chaire de Physique Mésoscopique
Michel Devoret
Année 2010, 11 mai - 22 juin

INTRODUCTION AU CALCUL QUANTIQUE
INTRODUCTION TO QUANTUM COMPUTATION

Sixième Leçon / *Sixth Lecture*

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or

<http://www.physinfo.fr/lectures.html>

[PDF FILES OF ALL PAST LECTURES ARE POSTED](#)

Questions, comments and corrections are welcome!

write to "phymeso@gmail.com"

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CONTENT OF THIS YEAR'S LECTURES

QUANTUM COMPUTATION FROM THE PERSPECTIVE OF MESOSCOPIC CIRCUITS

1. Introduction, c-bits versus q-bits
2. The Pauli matrices and quantum computation primitives
3. Stabilizer formalism for state representation
4. Clifford calculus
5. Algorithms
6. Quantum error correction

10-VL-3

CALENDAR OF SEMINARS

May 11: Cristian Urbina, (Quantronics group, SPEC-CEA Saclay)

Josephson effect in atomic contacts and carbon nanotubes

May 18: Benoît Douçot (LPTHE / Université Pierre et Marie Curie)

Towards the physical realization of topologically protected qubits

June 1: Takis Kontos (LPA / Ecole Normale Supérieure)

Points quantiques et ferromagnétisme

June 8: Cristiano Ciuti (MPQ, Université Paris - Diderot)

Ultrastrong coupling circuit QED : vacuum degeneracy and quantum phase transitions

June 15: Leo DiCarlo (Yale)

Preparation and measurement of tri-partite entanglement in a superconducting quantum circuit

June 22: Vladimir Manucharian (Yale)

The fluxonium circuit: an electrical dual of the Cooper-pair box?

10-VL-4

LECTURE IV : ERROR CORRECTION

Maintaining by an active feedback process,
not simply a single state, but a manifold of quantum states.

1. Classical error correction
2. A simplified example: 3-qubit code
3. Error processes
4. The 7-qubit code

10-VI-5

OUTLINE

1. Classical error correction
2. A simplified example: 3-qubit code
3. Error processes
4. The 7-qubit code

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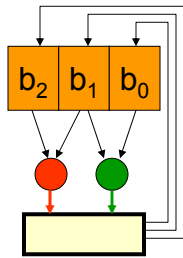
BASICS OF CLASSICAL ERROR-CORRECTION

Redundancy: $\begin{matrix} 0 \\ 1 \end{matrix} \left. \vphantom{\begin{matrix} 0 \\ 1 \end{matrix}} \right\} \text{replaced by: } \begin{matrix} 0 & 0 & 0 \\ 1 & 1 & 1 \end{matrix} \quad \text{"repetition code"}$

Possible 1-bit errors: $\begin{matrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{matrix}$

Probability: ε per unit time (2-bit errors: ε^2)

Feedback information on parity of first 2 bits and last 2 bits:

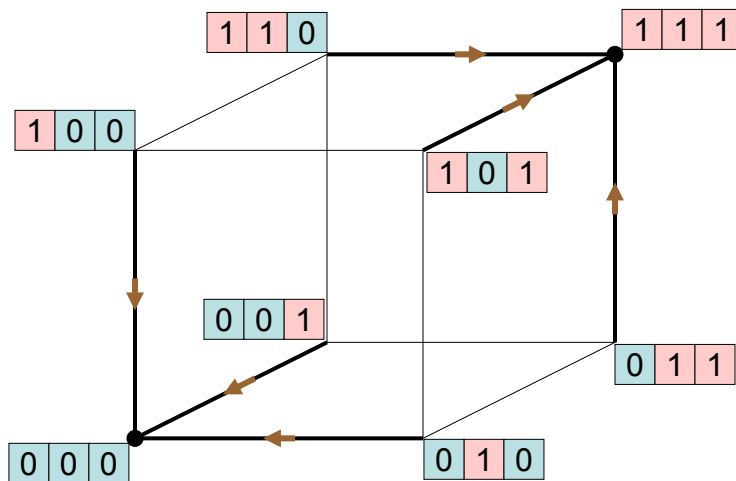


$$\begin{cases} e = b_2 \oplus b_1 \\ f = b_1 \oplus b_0 \end{cases}$$

$e \backslash f$	0	1
0	do nothing	$b_0 \rightarrow b_0 \oplus 1$
1	$b_2 \rightarrow b_2 \oplus 1$	$b_1 \rightarrow b_1 \oplus 1$

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PRINCIPLE OF CLASSICAL ERROR-CORRECTION



1 error is not lethal, system is kept within basin of attraction of protected state

10-VI-8

GENERAL PARITY ERROR BIT

$$N+1 \text{ bits} \quad \vec{x} = (x_N, x_{N-1}, \dots, x_2, x_1, x_0) \in \mathbb{B}^{N+1}$$

$$\text{constraint:} \quad e = \sum_{i=0}^N \oplus x_i \quad \leftarrow \quad \begin{array}{l} \text{Boolean sum,} \\ N \text{ free bits} \end{array}$$

↑
parity error bit, normal state is $e = 0$

If 1 or an odd number of errors occur, constraint is violated,
as shown by $e = 0 \rightarrow e = 1$.

It is possible to detect that an error has occurred,
but it is impossible to correct it, if nothing is added.

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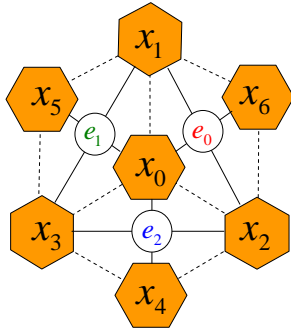
ERROR CORRECTION REQUIRES EXTRA BITS

Increase number of independent constraints on the bits until:

$$2 \text{ number of error bits} \geq \text{number of code bits} + 1$$

10-VI-9bis

HAMMING ERROR CORRECTING CODES



Example: 7 bits protected with 3 error bits, coding for 4 free bits ($2^3=7+1$)

Constraints:

$$x_0 \oplus x_1 \oplus x_2 \oplus x_6 = e_0 = 0$$

$$x_0 \oplus x_1 \oplus x_3 \oplus x_5 = e_1 = 0$$

$$x_0 \oplus x_2 \oplus x_3 \oplus x_4 = e_2 = 0$$

can be written as: $\mathbf{A}\vec{x} = \vec{e} = \vec{0}$

where:

$$\mathbf{A} = \begin{bmatrix} 1 & 0 & 0 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 0 & 1 \end{bmatrix}$$

$i \leftarrow$
 $\downarrow j$

After one error:

$$\mathbf{A}\vec{x}' = \vec{e} \neq \vec{0}$$

The error syndrome matrix \mathbf{A} is used to detect which error has occurred **and** to correct it.

$$x_i \rightarrow x_i \oplus \prod_{j=0}^2 (e_j \oplus \bar{A}_{ji})$$

Requires seven 3-way AND + linear gates

10-VI-10e

CORRECTING QUBIT ERRORS?!

- Necessary, no internal stabilizing dynamics as in c-bits.
- Cannot check errors directly: measurement destroys state.
- Qubit errors, unlike c-bit errors, occur continuously.
- Bit flips are not the only errors. Must correct phase flips.

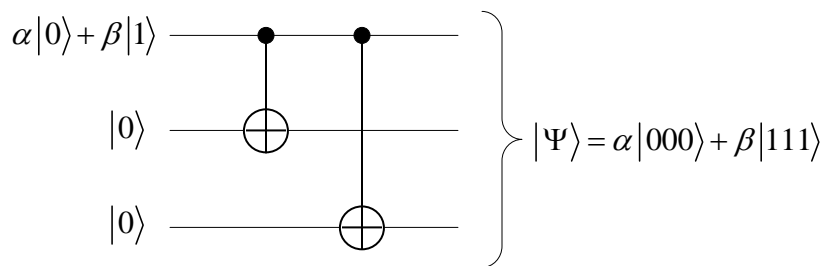
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OUTLINE

1. Classical error correction
2. A simplified example: 3-qubit code
3. Errors and error syndromes
4. The 7-qubit code

10-VI-5b

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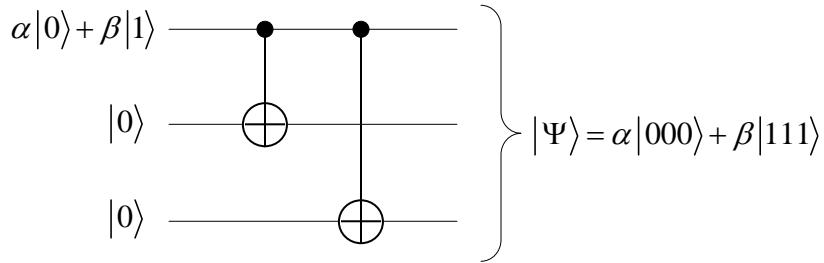
$$\begin{Bmatrix} +IZI \\ +IIZ \end{Bmatrix}$$

$$\begin{Bmatrix} +ZZI \\ +ZIZ \end{Bmatrix}$$

↑
stabilizer
of qubit
manifold

10-VI-12b

ENCODING



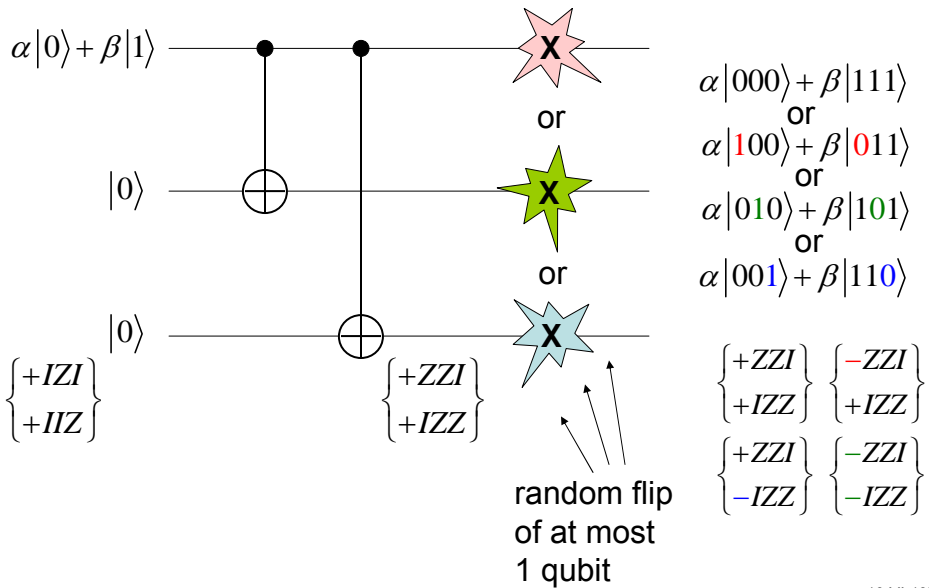
$$\begin{Bmatrix} +IZI \\ +IIZ \end{Bmatrix}$$

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↑
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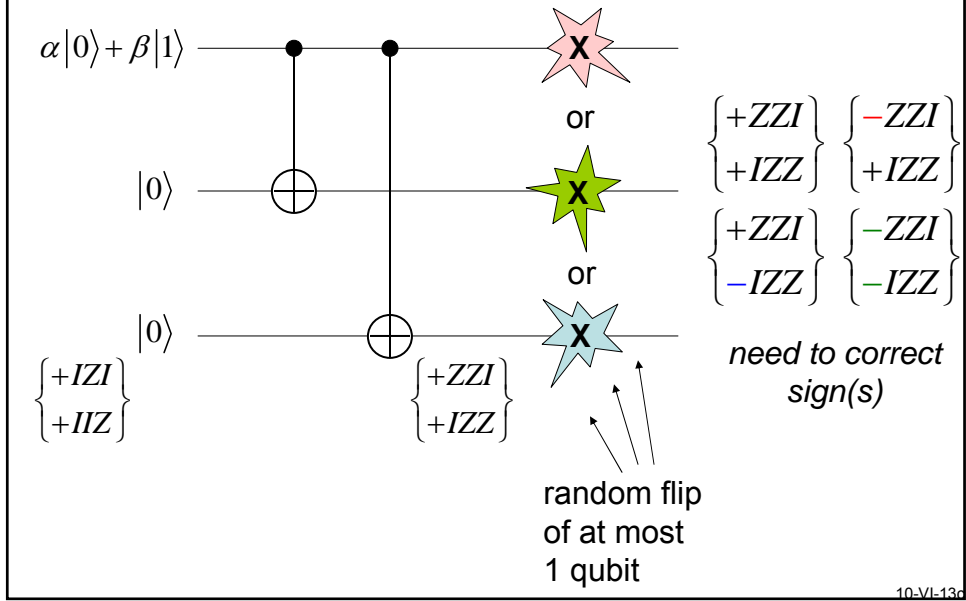
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BIT FLIP ERRORS



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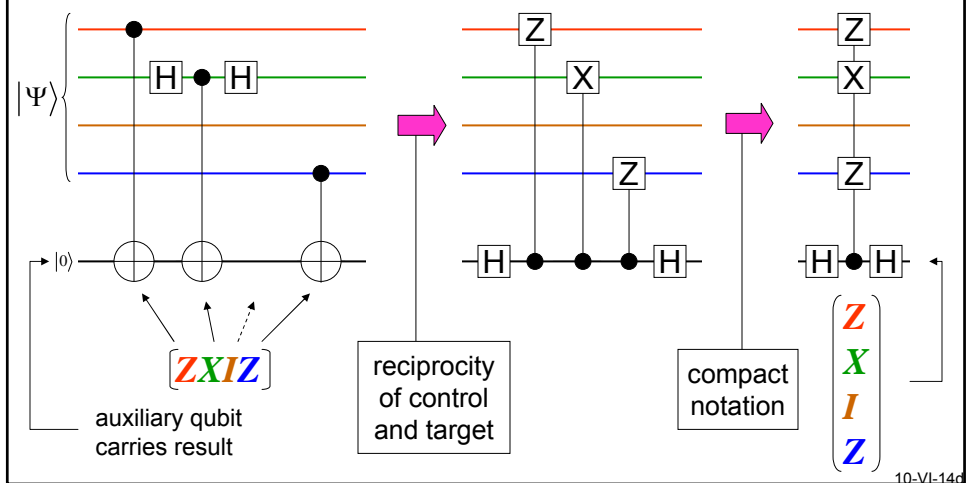
BIT FLIP ERRORS

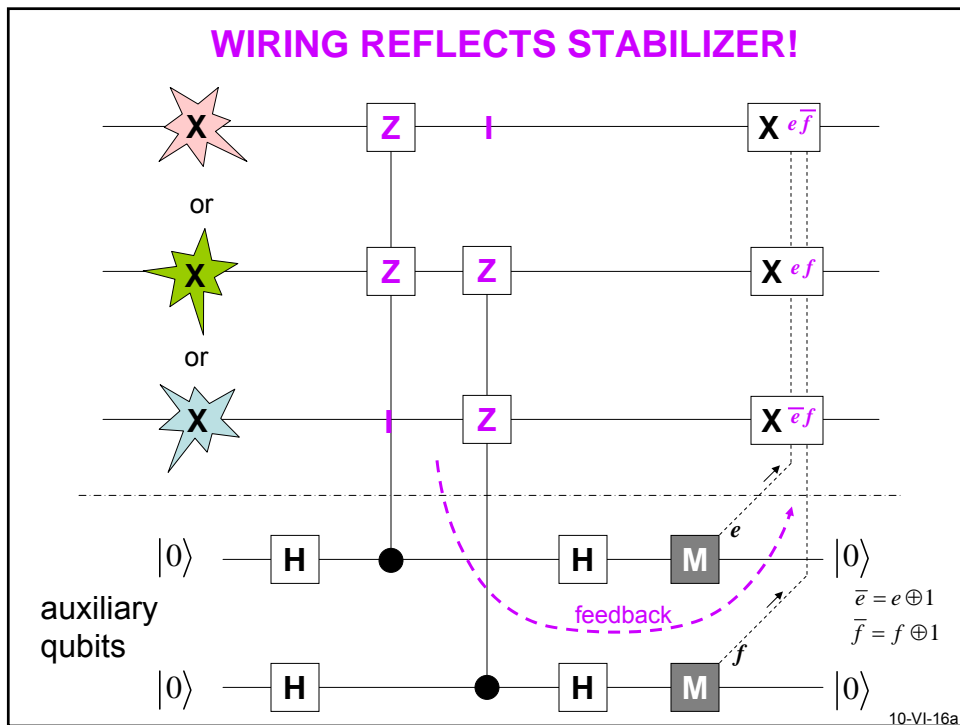
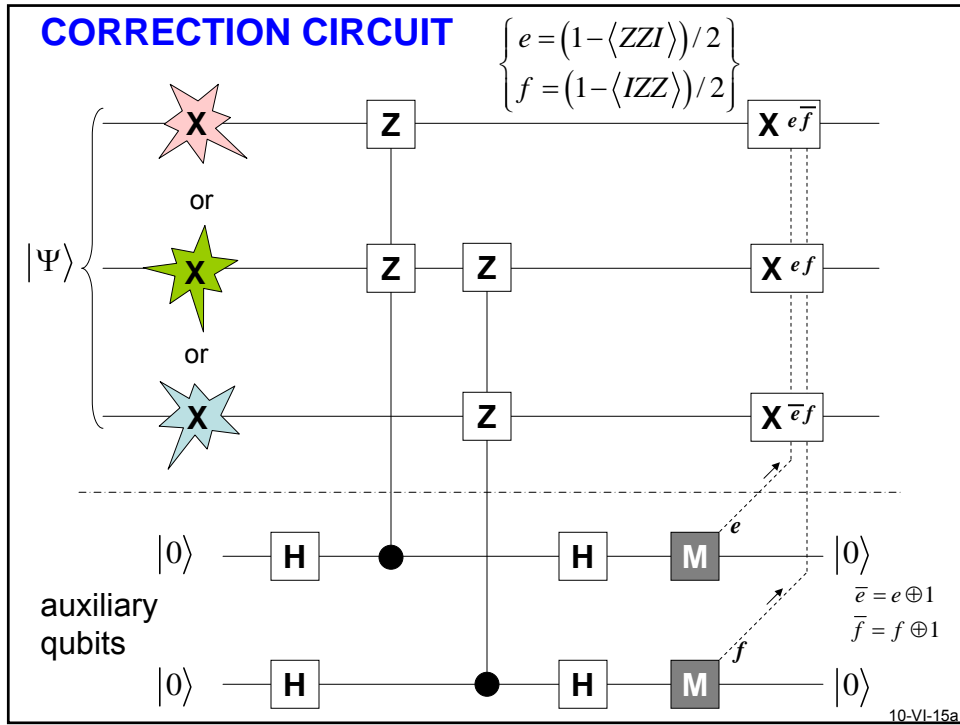


JOINT MEASUREMENTS

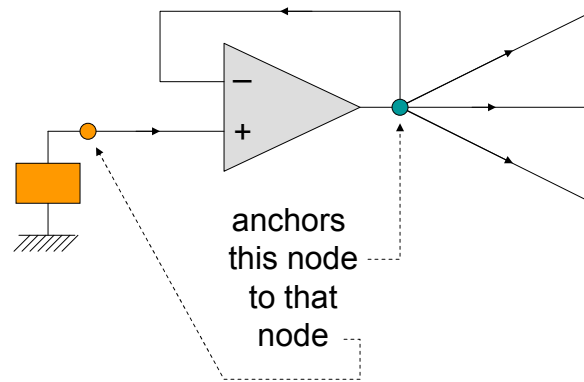
How do we measure joint qubit operators such as ZZI ?

Let us take a more representative case, for instance $ZXIZ$:





QUBIT CORRECTION CIRCUIT WORKS LIKE ORDINARY FEEDBACK



10-VI-16bis

POINTS FOR DISCUSSION

- 1) Correction protocol is discrete but error process is continuous.
How are "partial errors" dealt with?
- 2) Feedback goes through external circuitry.
Can feedback be purely internal to the system?
- 3) Error correction removes entropy from qubit.
Where is the entropy going?
- 4) Quantum error correction : feedback to a manifold, not a state.
What symmetry allows this manifold-preserving attractive dynamics?

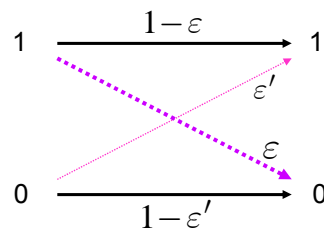
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OUTLINE

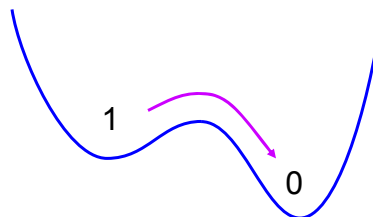
1. Classical error correction
2. A simplified example: 3-qubit code
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4. The 7-qubit code

10-VI-5c

EVEN CLASSICALLY, THERE IS MORE TO ERRORS THAN JUST RANDOM BIT FLIPS

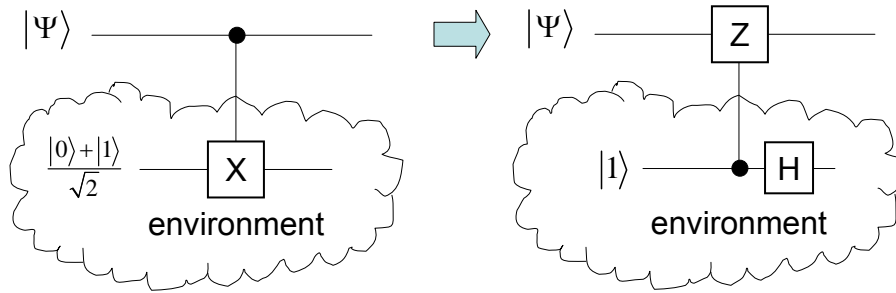


The two error transitions might not have the same probability



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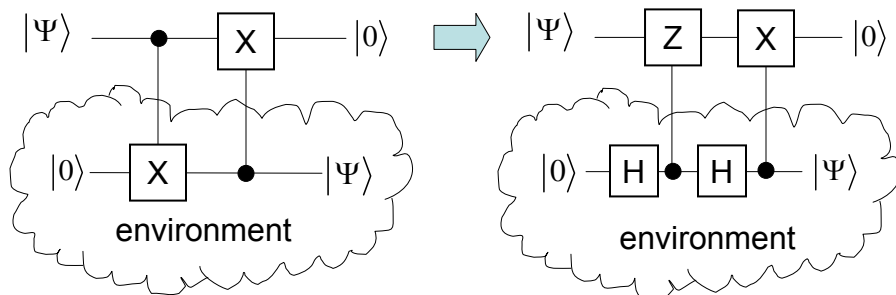
QUANTUM-MECHANICALLY, ERRORS CAN BE BIT FLIPS, PHASE FLIPS AND BIT-PHASE FLIPS.



Example of phase flip, leading to dephasing

10-VI-19a

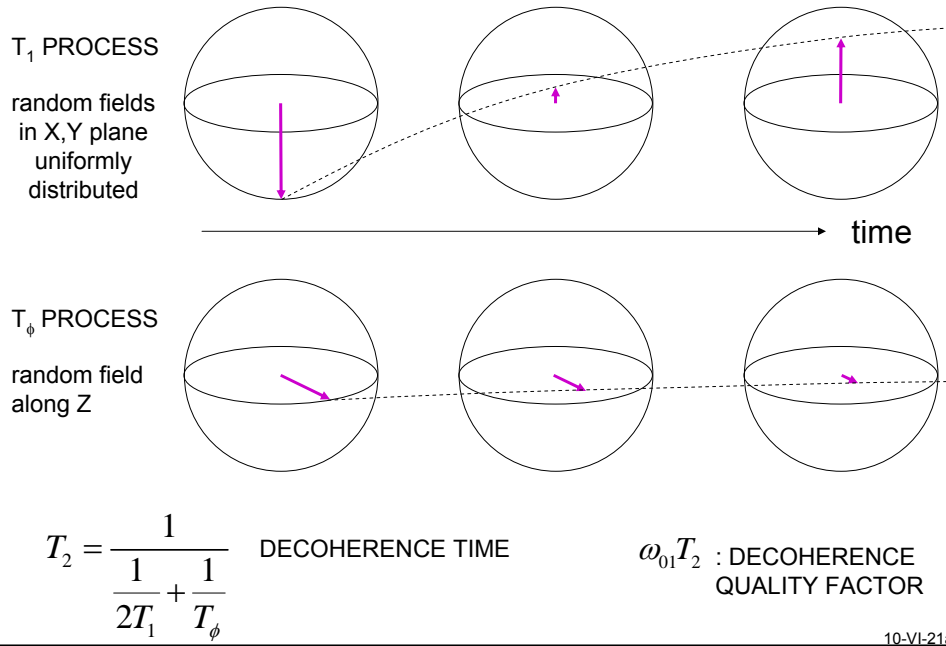
QUANTUM-MECHANICALLY, ERRORS CAN BE BIT FLIPS, PHASE FLIPS AND BIT-PHASE FLIPS (2)



Relaxation can be seen as a combination of phase and bit flips performed by a "cold" environment

10-VI-20a

LOSSES OF QUANTUM MEMORY



**JUST BY CORRECTING BIT FLIPS,
 PHASE FLIPS AND THE COMBINATION
 OF THESE TWO FLIPS,
 ANY TYPE OF ERROR CAN BE CORRECTED!**

Shor (1995), Steane (1996)

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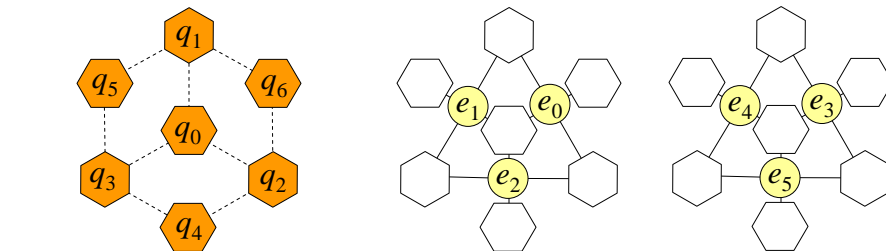
OUTLINE

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STABILIZER OF THE 7-QUBIT CODE

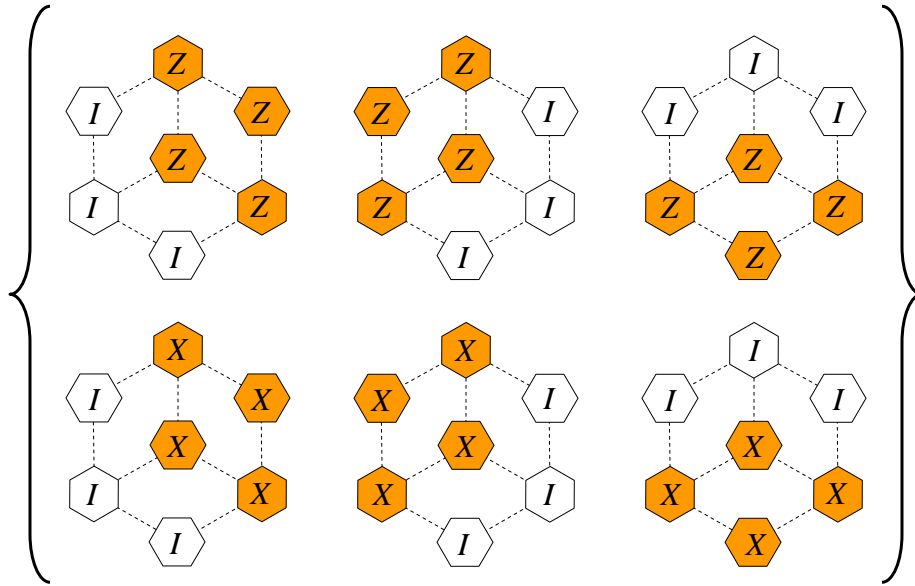
Steane (1996)



$$\mathbf{A} = \begin{matrix} & & & & i \leftarrow \\ \begin{matrix} q_i \\ \downarrow \\ \left[\begin{array}{ccccccc} 1 & 0 & 0 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 0 & 1 \end{array} \right] \\ \downarrow \\ e_j \end{matrix} & \rightarrow & \left\{ \begin{array}{l} +Z \quad I \quad I \quad I \quad Z \quad Z \quad Z \\ +I \quad Z \quad I \quad Z \quad I \quad Z \quad Z \\ +I \quad I \quad Z \quad Z \quad Z \quad I \quad Z \\ \dots \\ +X \quad I \quad I \quad I \quad X \quad X \quad X \\ +I \quad X \quad I \quad X \quad I \quad X \quad X \\ +I \quad I \quad X \quad X \quad X \quad I \quad X \end{array} \right\}
 \end{matrix}$$

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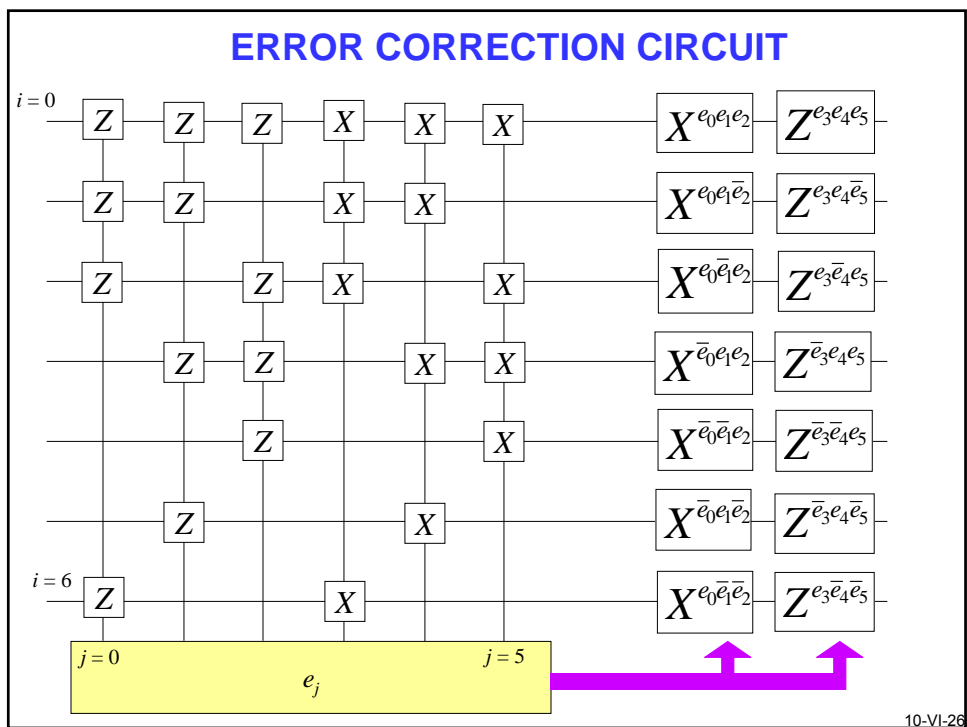
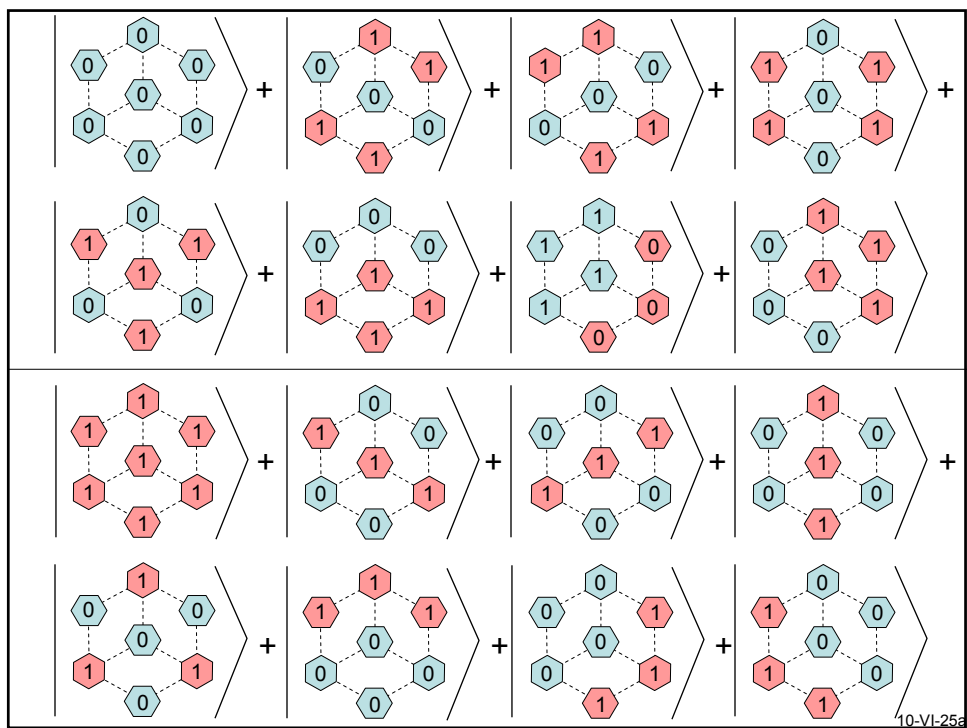
OTHER REPRESENTATION OF STEANE STABILIZER



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**NEXT SLIDE SHOWS THE TWO
WAVEFUNCTIONS OF STEANE CODE**

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WHY DOES IT WORK?

Steane's 7-qubit code:

6 generators in stabilizer (7 physical qubits – 1 logical qubit)

$2^6 = 64$ error syndromes > 7 qubits \times (3 errors/qubit) $+ 1 = 22$

Gottesman's 5-qubit code:

4 generators in stabilizer (5 physical qubits – 1 logical qubit)

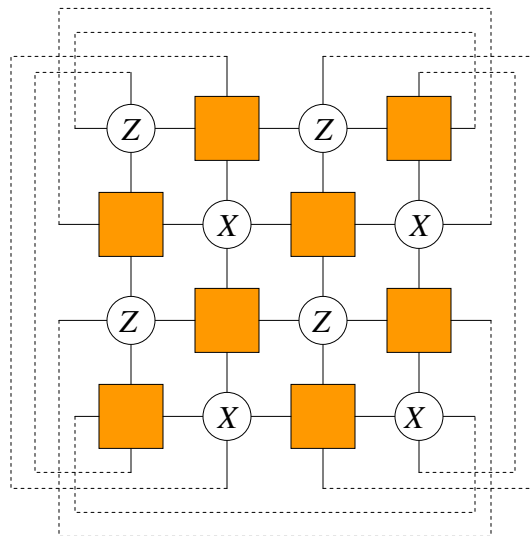
$2^4 = 16$ error syndromes $= 5$ qubits \times (3 errors/qubit) $+ 1 = 16$

↑
minimal but impractical

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TORIC CODE

Dennis, Kitaev, Landahl and Preskill (2001)



see B. Douçot's seminar this year

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ERRORS CAN ALSO BE CORRECTED BY CONTINUOUS MONITORING AND FEEDBACK

Ahn, Doherty & Landahl, Phys. Rev. A65, 042301 (2002)

NEXT YEAR: AMPLIFICATION AND FEEDBACK OF ENGINEERED QUANTUM SYSTEMS

10-VI-29

END OF 2010 COURSE

ACKNOWLEDGEMENTS

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